

$U_{2,4}$ —The only forbidden minor for binary matroids

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1 Background of Matroid Theory

- Definitions of matroids
- Representable matroids

2 The forbidden minor for binary matroids

- Binary matroids
- Theorem statement
- $U_{2,4}$ is a forbidden minor for binary matroids
- $U_{2,4}$ is the only forbidden minor for binary matroids

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Def: Independent sets

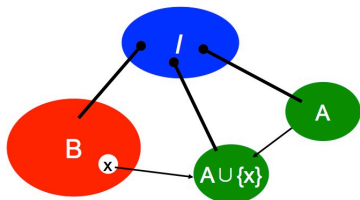
Definition

A matroid M is an ordered pair (E, \mathcal{I}) consisting of a finite set E and a collection \mathcal{I} of subsets of E having the following three properties:

- 1 $\phi \in \mathcal{I}$;
- 2 Hereditary:
If $I \in \mathcal{I}$ and $I' \subseteq I$, then $I' \in \mathcal{I}$;
- 3 Augmentation:
If I_1 and I_2 are in \mathcal{I} and $|I_1| < |I_2|$, then there is an element e of $I_2 - I_1$ such that $I_1 \cup e \in \mathcal{I}$.

Demo

Augmentation



A collection of subsets $\mathcal{B} \subseteq 2^E$ of a ground set E are the bases of a matroid if and only if:

- \mathcal{B} is non-empty;
- base exchange: If $B_1, B_2 \in \mathcal{B}$, for any $x \in (B_1 - B_2)$, there is an element $y \in (B_2 - B_1)$ such that $(B_1 - x) \cup y \in \mathcal{B}$.

Consider

$$A = \begin{array}{cccccc} & 1 & 2 & 3 & 4 & 5 & 6 \\ \left[\begin{array}{cccccc} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \end{array} \right] \end{array}$$

Bases: $\{1,2,3\}, \{1,2,4\},$
 $\{1,2,5\}, \{1,3,4\}, \{1,4,5\},$
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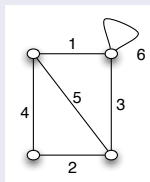
Def: Circuits

The circuits $\mathcal{C}(M)$ of a matroid are the minimal dependent sets.

Property

- If any one element is deleted from a circuit, it will become an independent set.
- Independent sets $\mathcal{I}(M)$ does not contain any element in $\mathcal{C}(M)$.
- Size 1 circuit is a loop
- Size 2 circuit contains parallel elements
- Simple matroids: no loops or parallel elements

Consider the graph

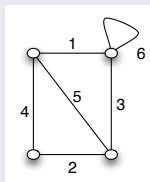


Circuits: $\{1,2,3,4\}$, $\{1,3,5\}$,
 $\{2,4,5\}$, $\{6\}$

A collection of sets \mathcal{C} is the collection of circuits of a matroid if and only if:

- $\emptyset \notin \mathcal{C}$
- if $C_1, C_2 \in \mathcal{C}$ with $C_1 \subseteq C_2$ then $C_1 = C_2$
- if $C_1, C_2 \in \mathcal{C}$, $C_1 \neq C_2$ and $e \in C_1 \cap C_2$ then there is some $C_3 \in \mathcal{C}$ with $C_3 \subseteq (C_1 \cup C_2) - e$.

Consider the graph



Circuits: $\{1,2,3,4\}$, $\{1,3,5\}$,
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Def: Rank function

For any base of a matroid,

$$r(B) = r(M).$$

Formally, a rank function is

$r : 2^E \rightarrow \mathbb{N} \cup \{0\}$, for which

- For $X \subseteq E$, $0 \leq r(X) \leq |X|$
- If $X \subseteq Y \subseteq E$, $r(X) \leq r(Y)$
- $\forall X, Y \subseteq E$, $r(X \cap Y) + r(X \cup Y) \leq r(X) + r(Y)$

Consider

$$A = \begin{array}{cccccc} & 1 & 2 & 3 & 4 & 5 & 6 \\ \left[\begin{array}{cccccc} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \end{array} \right] \end{array}$$

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Rank: $r(M) = r(B) = 3$

Given a rank function of a matroid,
the closure operation $\text{cl} : 2^E \rightarrow 2^E$ is
defined as

$$\text{cl } X = \{x \in E \mid r(X \cup x) = r(X)\}.$$

Closure is a span of a subspace.

$$\text{cl } E = E;$$

Independent Set:

$$\mathcal{I} = \{X \subseteq E \mid x \notin \text{cl}(X - x) \forall x \in X\}$$

A subset X of $E(M)$ is said to
be a flat if $\text{cl } X = X$.

$r(M) - 1$ flat is a hyperplane.

Consider

$$A = \begin{matrix} & 1 & 2 & 3 & 4 & 5 & 6 \\ \begin{bmatrix} 1 & 0 & 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \end{bmatrix} \end{matrix}$$

$\text{cl } \{1,2,6\} = \{1,2,6\} \rightarrow$ flat,
hyperplane

$\text{cl } \{1,3\} = \{1,3,5,6\}$

A function $\text{cl} : 2^E \rightarrow 2^E$ is closure for a matroid $M = (E, \mathcal{I})$ iff

- If $X \subseteq E$ then $X \subseteq \text{cl}X$
- If $X \subseteq Y \subseteq E$, then $\text{cl}X \subseteq \text{cl}Y$
- If $X \subseteq E$ then $\text{cl}(\text{cl}(X)) = \text{cl}X$
- If $X \subseteq E$ and $x \in E$ and $y \in \text{cl}(X \cup x) - \text{cl}(X)$ then $x \in \text{cl}(X \cup y)$

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$$\text{cl} \{1,2\} = \{1,2,6\} \rightarrow \text{flat}$$

$$\text{cl} \{1,3\} = \{1,3,5,6\}$$

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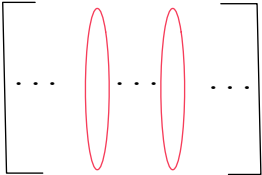
Representable Matroids

Definition

A matroid M with ground set S of size $|S| = N$ and rank $r_M = r$ is representable over a field \mathbb{F} if there exists a matrix $A \in \mathbb{F}^{r \times N}$ such that for each independent set $I \in \mathcal{I}$ the corresponding columns in A , viewed as vectors in \mathbb{F}^r , are linearly independent.

$$S_1, \dots, S_i, \dots, S_j, \dots, S_N \xleftrightarrow{\text{Mapping}} 1 \ 2 \ \dots \ i \ \dots \ j \ \dots \ N$$

$r(S_i, S_j) = \text{rank}(i, j\text{-th columns})$



The diagram shows a large square matrix with square brackets on the left and right sides. Inside the matrix, there are three horizontal ellipses representing rows. The first and third rows are partially visible. Two vertical red ellipses are drawn around the second and fourth columns of the matrix, highlighting them.

- Representable matroids usually can be characterized by forbidden minors (cannot contain such minors).

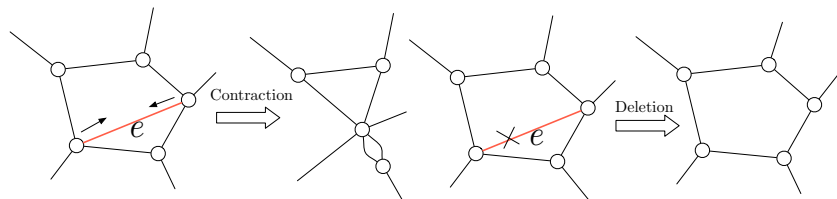
Definition

If M is a matroid on S and $T \subseteq S$, a matroid M' on T is called a *minor* of M if M' is obtained by any combination of deletion (\setminus) and contraction ($/$) of M .

Contraction and Deletion

Let M/T denote the matroid obtained by contraction of M on $T \subset S$, and let $M \setminus T$ denote the matroid obtained by deletion from M of $T \subset S$. Then, $\forall X \subseteq S - T$

$$\begin{aligned} r_{M/T}(X) &= r_M(X \cup T) - r_M(T) \\ r_{M \setminus T}(X) &= r_M(X) \end{aligned} \tag{1}$$



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Binary Matroids

- $M = M(A)$, where A is a binary matrix.
- Unique representation under equivalence: $A = [I_r | D]$, where I_r is identity matrix representing a basis B , $D_{e,f} = 1$ or 0 , if $B - e + f$ is or is not a basis.
- Example: selected basis $\{a, b, c\}$, other bases $\{b, c, d\}$, $\{a, b, d\}$, $\{b, c, e\}$, $\{a, c, e\}$

$$\begin{array}{c} a \\ b \\ c \end{array} \begin{bmatrix} & a & b & c & d & e \\ 1 & 0 & 0 & 1 & 1 \\ 0 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 \end{bmatrix}$$

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Theorem (Tutte 1958)

A matroid is binary if and only if it has no $U_{2,4}$ -minor.

$U_{2,4}$: rank=2, for any 2 or more of the total 4 elements.

Proof points:

- 1 $U_{2,4}$ is not binary representable and so its extensions. Hence, $U_{2,4}$ is a forbidden minor;
- 2 $U_{2,4}$ is the only such forbidden minor: suppose exists another forbidden minor, equivalent to $U_{2,4}$

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Point 1: $U_{2,4}$ is a forbidden minor

- $U_{2,4}$ is not binary representable due to limit of numbers of binary

vectors: $A = \begin{matrix} & 1 & 2 & 3 & 4 \\ \begin{bmatrix} 1 & 0 & 1 & 1 \\ 0 & 1 & 1 & ? \end{bmatrix} \end{matrix}$

- Extensions of a not- \mathbb{F}_q -representable matroid will not be \mathbb{F}_q representable: Suppose a matroid is \mathbb{F}_q -representable by a matrix, deletion/contraction on associated columns in the matrix will give representation for its minors, obtained by some deletion/contraction.
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Point 2.0: Suppose existence of another forbidden minor

- Forbidden minors: simple, not binary, but minors are binary.
- Suppose exists another forbidden minor M .
- M has no $|E(M)| - 1$ - or $|E(M)| - 2$ -element hyperplane;

- if has $|E(M)| - 1$ -element hyperplane, representation for M by adding column to the representation matrix $[I_{r-1}|D]$ of the hyperplane:

$$\begin{bmatrix} I_{r-1} & \mathbb{O} & D \\ 0 & 1 & \mathbb{O} \end{bmatrix}$$

- if has $|E(M)| - 2$ -element hyperplane, representation will be:

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Point 2.1: Representation and M'

- For any distinct $x, y \in E(M)$, $r(M \setminus \{x, y\}) = r(M)$.
- Suppose $[I_r|D] = M \setminus \{x, y\}$ over binary. $[I_r|D|v_x] = M \setminus \{y\}$ and $[I_r|D|v_y] = M \setminus \{x\}$. Let $M' = [I_r|D|v_x|v_y]$.
- $M \setminus x = M' \setminus x$ and $M \setminus y = M' \setminus y$, but $M \neq M'$.
- \exists a set s.t. independent in one but dependent in the other.
- Let Z be such a minimal set.
 $Z \in \mathcal{I}(M_I), Z \in \mathcal{C}(M_C), \{M_I, M_C\} = \{M, M'\}$. (if $Z \notin \mathcal{C}(M_C)$, \exists some elements to delete and keep the dependence, contradicting with the minimality)
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Point 2.2: $Z = \{x, y\}$ & $r_{M_I}(E) = 2$

- Want to show: $Z = \{x, y\}$ & $r_{M_I}(E) = 2$
- Consequences of: if $Z \subseteq J, J \in \mathcal{I}(M_I)$, then $J = \{x, y\}$.
- Suppose $G = J - \{x, y\} \neq \emptyset$. $G \in \mathcal{I}(M_I \setminus \{x, y\})$ and $G \in \mathcal{I}(M_C \setminus \{x, y\})$.
- $N_I = M_I/G, N_C = M_C/G$, we have $r_{N_I}(E \setminus G) = r_{N_C}(E \setminus G) = r_{M_I}(E) - r_{M_I}(G) = r_{M_C}(E) - r_{M_C}(G)$, but $N_I \neq N_C$ (after contraction, x, y are dependent in N_C , while independent in N_I).
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- $N_I = M_I/G, N_C = M_C/G$, we have $r_{N_I}(E \setminus G) = r_{N_C}(E \setminus G) = r_{M_I}(E) - r_{M_I}(G) = r_{M_C}(E) - r_{M_C}(G)$, but $N_I \neq N_C$ (after contraction, x, y are dependent in N_C , while independent in N_I).
- However, still have $N_I \setminus x = N_C \setminus x$ and $N_I \setminus y = N_C \setminus y$.

Point 2.2: $Z = \{x, y\}$ & $r_{M_I}(E) = 2$

- Consider a basis B for $N_I \setminus \{x, y\}$, which equals $N_C \setminus \{x, y\}$. B is either a basis of both N_I, N_C or of neither. (reason: N_I, N_C same rank)
- If B is a basis of neither, $r_{N_I}(E \setminus G \setminus \{x, y\}) < r_{N_I}(E \setminus G)$, contradicting with the fact that no $|E(M)| - 1$ - or $|E(M)| - 2$ -element hyperplane.
- If add back $r_{N_I}(G)$ to both sides, get $r_{M_I}(E \setminus \{x, y\}) < r_{N_I}(E)$

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Point 2.3: $M = U_{2,4}$

- B is a basis of both N_I, N_C , both are binary.
- $N_I \setminus x = N_C \setminus x = [I_{r'} | D' | v'_y]$ and $N_I \setminus y = N_C \setminus y = [I_{r'} | D' | v'_x]$,
 $[I_{r'} | D' | v'_x | v'_y] = N_I = N_C$ due to uniqueness. Contradiction.
- $\{x, y\} \subseteq Z \subseteq J = \{x, y\}$. A basis $J \supseteq Z$ is $\{x, y\}$, we have
 $r_{M_I}(E) = 2$.
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- Consider every possible x, y , M , either M_I, M_C , has rank 2 and $\{x, y\}$
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Thank you!